



# Matching with Transfers under Distributional Constraints

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## Introduction and Motivation

- Many two-sided matching markets with **transferable utilities**, e.g., labor or rental housing markets, are subject to **distributional constraints**
- Prior work on matching under constraints has mainly focused on the non-transferable utility setting
- We study conditions on the constraint structure under which **equilibria exist** in the transferable utility setting



## Definitions

An **arrangement** is specified by an assignment  $\mathbf{X}$  and salaries  $\{s_f\}_{f \in F}$ , where  $F$  is the set of firms

### Key Properties

**1. Feasibility:** An arrangement is feasible if the assignment  $\mathbf{X}$  satisfies the distributional constraints

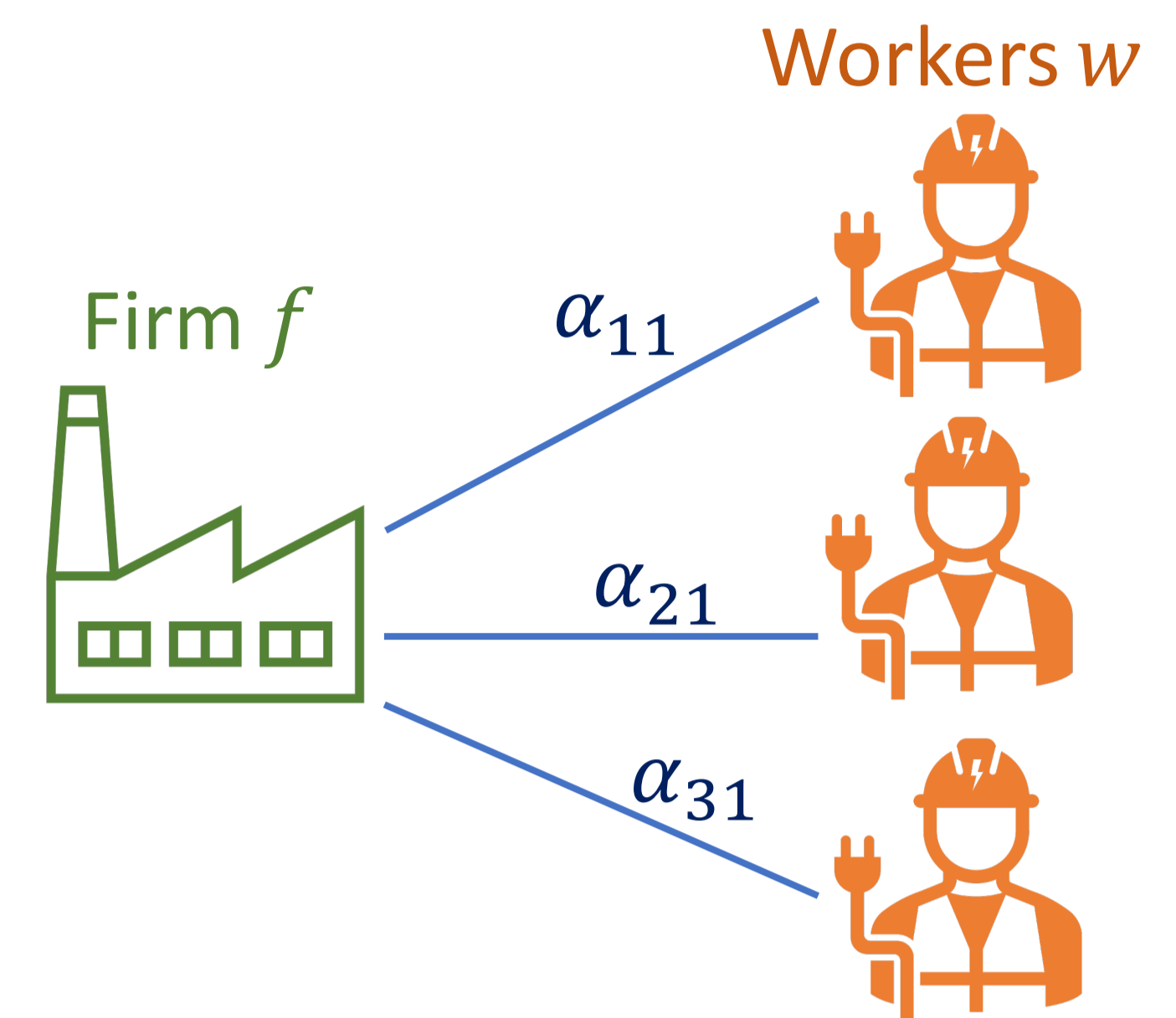
**2. Stability:** No firm and *feasible* group of workers can form a coalition, and all become better off by deviating

**3. Efficiency:** An assignment is *efficient* if it maximizes the total match value among all *feasible* assignments

**Theorem (Efficiency of Stable Arrangements):** If a feasible arrangement is stable, then the corresponding assignment is efficient.

## Model

### Total Match Value



$a_{wf}$ : Value of firm  $f$  for worker  $w$

$c_{wf}$ : Value of worker  $w$  for firm  $f$

$\alpha_{wf} = a_{wf} + c_{wf}$ : Total Match Value of worker-firm pair  $w-f$

### Payoffs

**Worker Payoff** on matching with  $f$ :

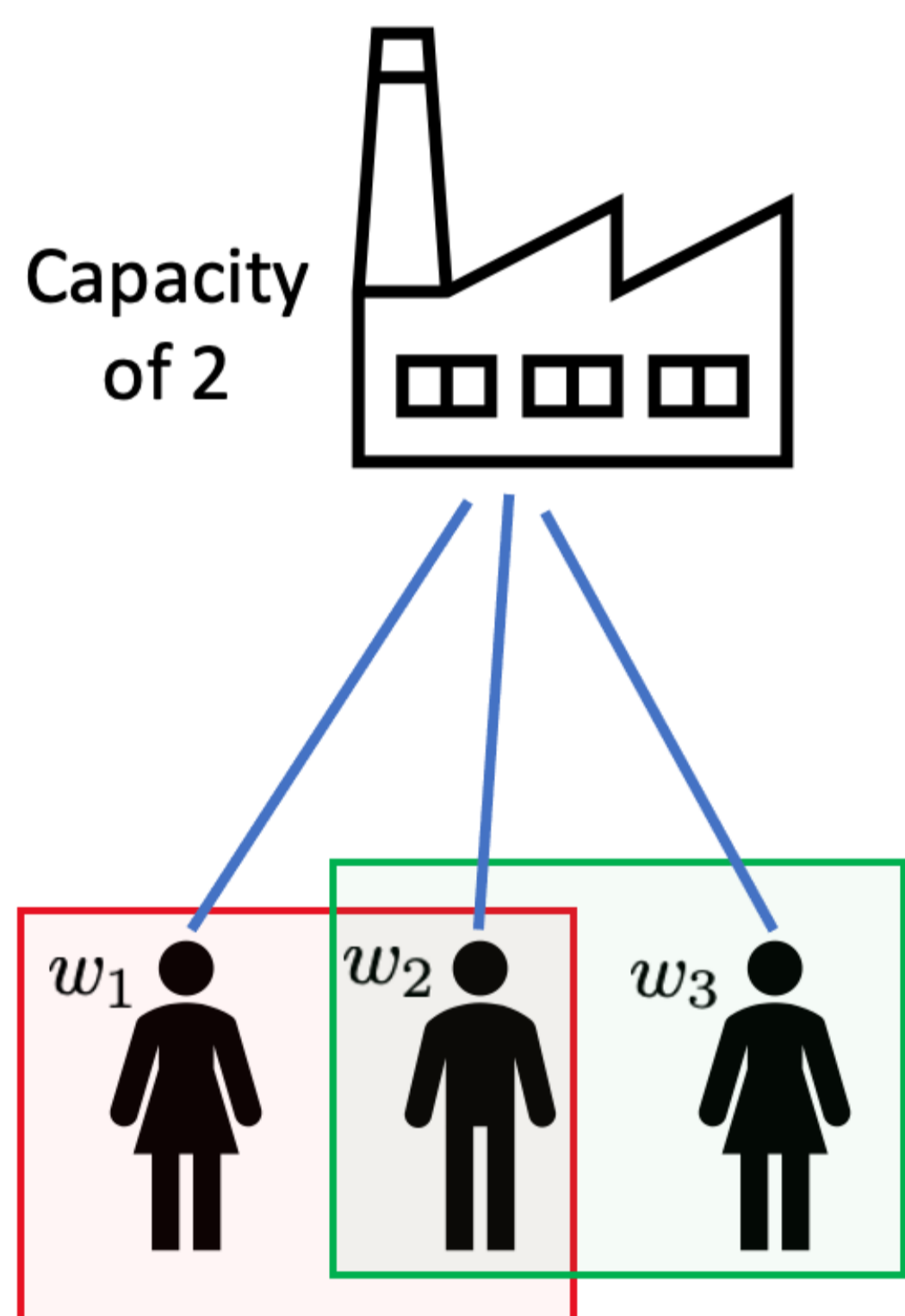
$$u_w = a_{wf} + s_{wf}$$

**Firm Payoff** on matching with feasible set of workers  $D \in T_f$ :

$$v_f = \sum_{w \in D} (c_{wf} - s_{wf})$$

## Stable Arrangements for One Firm Setting

### Example where Substitutes Condition is Violated



Value of Worker for Firm	Value of Firm for Worker
$a_{11} = -0.5$	$c_{11} = 1.5$
$a_{21} = -0.5$	$c_{21} = 2.5$
$a_{31} = -0.5$	$c_{31} = 1.5$

Salary Profile 1	Salary Profile 2
$s_{11} = 0.5$	$s'_{11} = 1.1$
$s_{21} = 1$	$s'_{21} = 1$
$s_{31} = 0.5$	$s'_{31} = 0.5$

$$x_{w_1, f_1} + x_{w_2, f_1} \leq 1$$

$$x_{w_2, f_1} + x_{w_3, f_1} \leq 1$$

Increasing the salary of worker one decreases firm's demand for worker three

**Substitutes condition:** If the wage for one worker is increased, the demand for all other workers should weakly increase

### Theorem (Existence of Stable Arrangement for One Firm):

In a one firm setting, stable arrangements exist irrespective of the nature of the constraint structure or agent's preferences.

**Takeaway: Stable arrangements exist even when substitutes condition is violated**

## Future Work

- Extending linear programming approach to a broader range of firm preferences
- Extending results to setting when preferences are not quasi-linear
- Developing mechanisms that are incentive compatible in addition to preserving stability

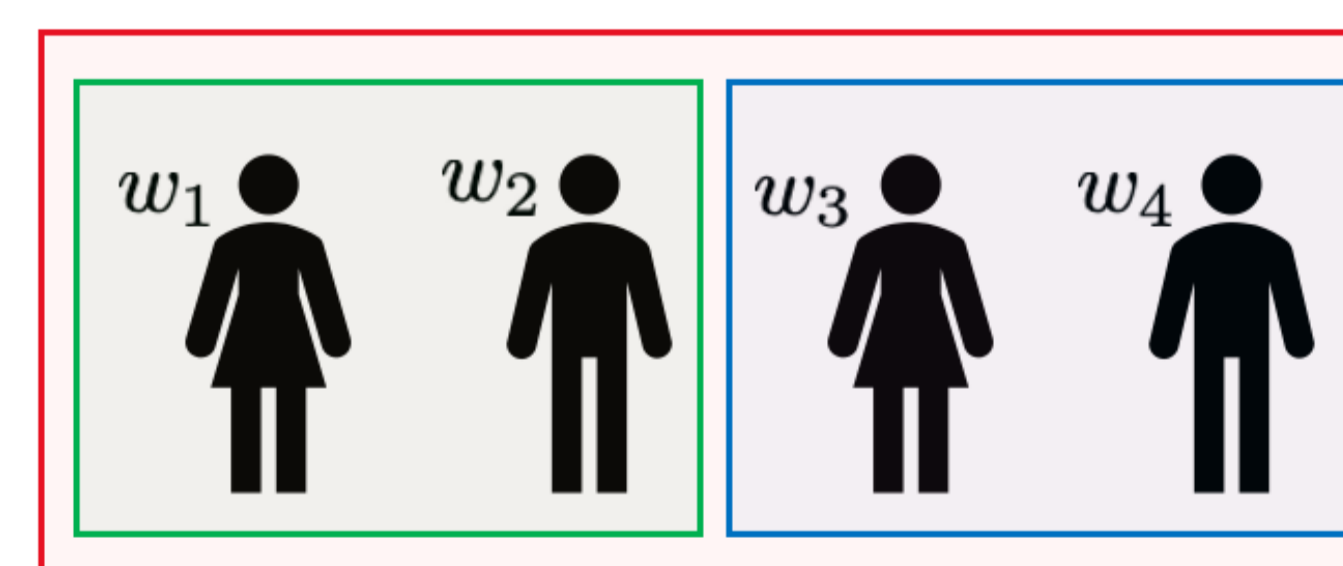
## Stable Arrangements in Multiple Firm Setting

**Proposition:** Under general constraints, stable arrangements may not exist, even if firms have linear preferences

$$\begin{aligned} \max_{\mathbf{X} \in \mathbb{R}^{|W| \times |F|}} \quad & U(\mathbf{X}) = \sum_{w \in W} \sum_{f \in F} \alpha_{w,f} x_{w,f}, && \text{Total Match Value} \\ \text{s.t.} \quad & \sum_{f \in F} x_{w,f} \leq 1, \quad \forall w \in W, && \text{Worker Allocation Constraints} \\ & x_{w,f} \geq 0, \quad \forall w \in W, f \in F, \\ & \sum_{w \in D} x_{w,f} \leq \bar{\lambda}_D, \quad \forall D \in \mathcal{H}_f, f \in F. && \text{Distributional Constraints} \end{aligned}$$

**Theorem (Linear Programming and Stable Arrangements):** If the distributional constraints form a polymatroid (or hierarchy), then there exists a stable arrangement, and it can be computed in polynomial time using linear programming.

### Hierarchy/Polymatroid



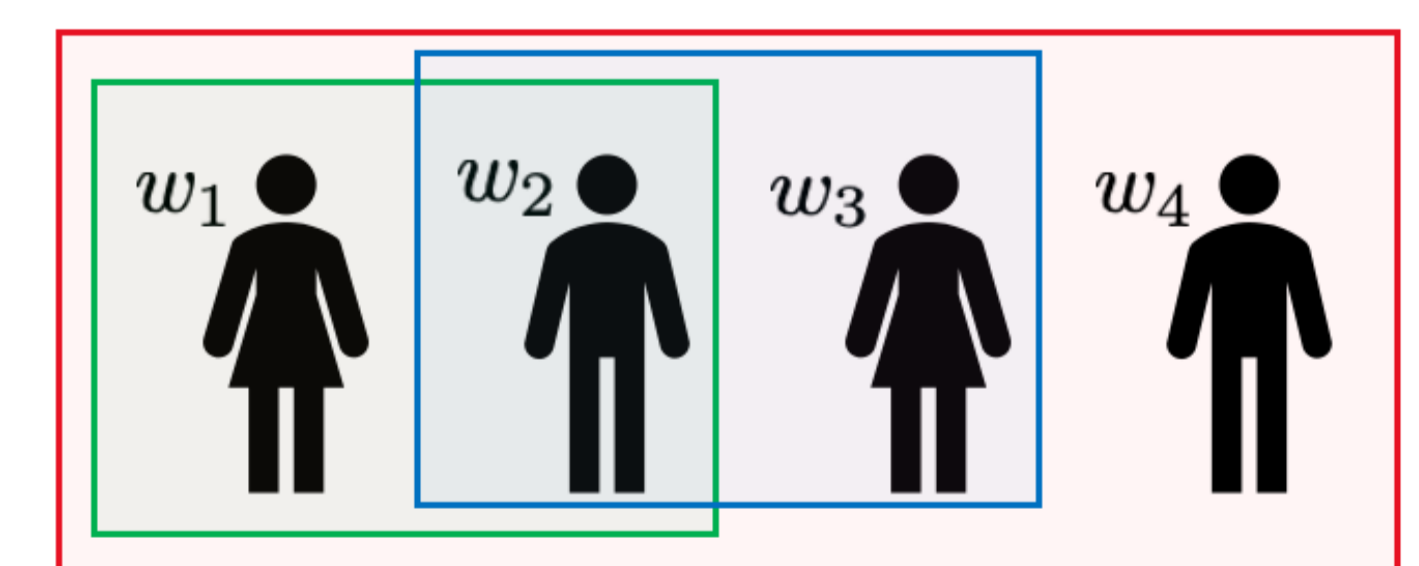
Constraints

$$x_{w_1, f_1} + x_{w_2, f_1} + x_{w_3, f_1} + x_{w_4, f_1} \leq 2$$

$$x_{w_1, f_1} + x_{w_2, f_1} \leq 1$$

$$x_{w_3, f_1} + x_{w_4, f_1} \leq 1$$

### Not a Hierarchy/Polymatroid



Constraints

$$x_{w_1, f_1} + x_{w_2, f_1} + x_{w_3, f_1} + x_{w_4, f_1} \leq 2$$

$$x_{w_1, f_1} + x_{w_2, f_1} \leq 1$$

$$x_{w_2, f_1} + x_{w_3, f_1} \leq 1$$

**Hierarchy:** For all sets of workers  $S, S' \in \mathcal{H}_f$ , either  $S \subseteq S'$ ,  $S' \subseteq S$  or  $S \cap S' = \emptyset$

Linear Programming approach can be extended to **lower bound quotas** under a slight modification to the stability notion